



General and Skewed General Variable Neighborhood Search Approaches for Integrated Planning and Scheduling

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1. Introduction

- Planning and scheduling processes are strongly interrelated, but they are usually addressed separately.
- To achieve **global optimal solutions**, the related optimization problems must be addressed all together in an **integrated way** [1].
- Here, we consider the Integrated Planning and Scheduling Problem (IPSP), as it was firstly proposed in [2].

2. Objectives

- New **heuristics and metaheuristic methods** to tackle the IPSP on parallel and identical machines [2];
- Innovative neighborhood structures** specially designed for the IPSP to be used within **Variable Neighborhood Search** algorithm (VNS) and its variants **General VNS (GVNS)** and **Skewed General VNS (SGVNS)**.

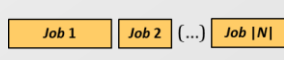
3. Problem Definition

The entire time horizon (T) is divided into τ **time periods** of length equal to $P \in \mathbb{N}$. There are $|M|$ **parallel and identical machines**, where the **set of jobs** N should be processed. Each **job** $j \in N$ has five attributes: **processing time** – $p_j \in \mathbb{N}$; **release date** – $r_j \in T$; **due date** – $d_j \in T$; **penalty factors** incurred when a job is completed before and after its due date – $e_j, l_j \in \mathbb{N}$, respectively.

Machines available time in each t period (bins)



Jobs processing time (objects/packages)



The objective of this IPSP is to **minimize the sum of all penalties** (w_j^t), where t represents the period in which job j will be processed:

$$w_j^t = e_j \times \underbrace{\max\{0, d_j - t\}}_{\text{earliness}} + l_j \times \underbrace{\max\{0, t - d_j\}}_{\text{tardiness}}$$

4. Neighborhood Structures

Table 1. Set of 12 jobs and attributes (example).

(job) j	1	2	3	4	5	6	8	9	10	11	12
r_j	1	1	1	1	1	1	1	2	3	1	1
d_j	1	1	1	1	1	2	3	3	3	1	3
p_j	3	2	2	4	4	2	2	1	4	2	3
e_j	0	0	0	0	0	10	4	7	0	0	5
l_j	10	5	12	10	9	10	8	12	10	3	8

$N = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12\}$,
 $M = \{1, 2\}$, $P = 8$, $\tau = 3$

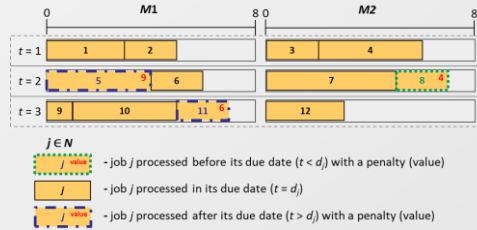


Fig. 1. A feasible solution S .

5. Computational Experiments

Experiments were performed on benchmark instances divided into three sets: **small jobs** (set A), **large jobs** (set C), and **half-half** (set B). Five runs for each instance (time limit of 5 seconds). The **VNS**, **GVNS** and **SGVNS** results were compared with the best results from **Exact Methods (EM)** [3].

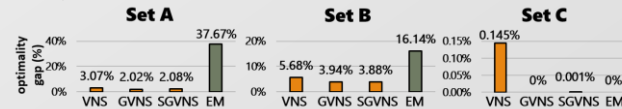


Fig. 8. Average **optimality gap** (in %).

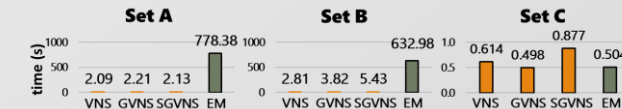


Fig. 9. Average **computational time** (in seconds).

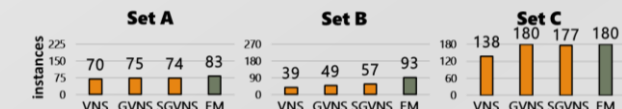


Fig. 10. Number of instances solved up to **optimality**.



Fig. 2. \mathcal{N}_1 – Insertion of one job in another period. Solution $S' \in \mathcal{N}_1(S)$.



Fig. 3. \mathcal{N}_2 – Exchange between two jobs from different periods. Solution $S' \in \mathcal{N}_2(S)$.

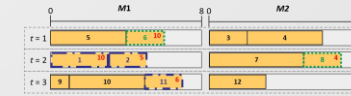


Fig. 4. \mathcal{N}_3 – Exchange between two different bins of different periods. Solution $S' \in \mathcal{N}_3(S)$.



Fig. 5. \mathcal{N}_4 – Exchange between two jobs from the same period, but from different machines, and an insertion of another job in the considered period. Solution $S' \in \mathcal{N}_4(S)$.



Fig. 6. \mathcal{N}_5 – Exchange between two jobs from necessarily adjacent periods. Solution $S' \in \mathcal{N}_5(S)$.



Fig. 7. \mathcal{N}_6 – Exchange between two jobs, from the same period, but from different machines. Solution $S' \in \mathcal{N}_6(S)$.

6. Conclusions

The proposed approaches are able to obtain solutions very close to the optimal solutions (**low values of the optimality gap**). The major core of these approaches is the **very short time to achieve good quality solutions**, that are quite near the optimal solutions.

7. References

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